

ELGAMAL DIGITAL SIGNATURES

Elgamal signature scheme involves the use of private key for encryption and public key for decryption

The global elements of Elgamal digital signature are prime number q and a , which is the primitive root of q .

1. Global Public key Components

- q - prime no.
 - a - primitive root of q
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2. User A signs a message M to B by computing

- Generate a random integer X_A , such that $1 < X_A < q-1$
- Compute $Y_A = a^{X_A} \text{ mod } q$
- A's Private key is X_A
- A's Public key is Y_A

To sign a message M , user A first computes the hash $m = H(M)$, such that m is an integer in the range $0 \leq m \leq (q-1)$

3. User A generates the digital signature

- Choose a random integer K , such that $1 \leq K \leq (q-1)$ and $\text{gcd}(K, q-1) = 1$. That is, K is relatively prime to $q-1$.
- Compute, $S_1 = a^K \text{ mod } q$
- Compute $K^{-1} \text{ mod } q-1$
- Compute, $S_2 = K^{-1}(m - X_A S_1) \text{ mod } (q-1)$
- The signature consists of a pair (S_1, S_2)

2. User B verifies the Signature

$$V_1 = a^m \text{ mod } q$$

$$V_2 = (Y_A)^{S_1} (S_1)^{S_2} \text{ mod } q$$

The signature is valid if $V_1 = V_2$.

Example I

Global Element

$q=19$ and $a=10$

Alice computes the private and public key

- Alice computes her key:
 - Alice chooses Private key, $X_A=16$
 - Computes Public Key, $Y_A=10^{16} \bmod 19 = 4$
- Alice signs message with hash $m=14$
 - Alice chooses $K=5$ which is relatively prime to $q-1=18$
 - Compute $S_1 = 10^5 \bmod 19 = 3$
 - Compute $K^{-1} \bmod (q-1) = 5^{-1} \bmod 18 = 11$
 - Compute $S_2 = 11(14-16*3) \bmod 18 = -374 \bmod 18 = 4 \{-374 \bmod 18 = 18 - 374 \% 18\}$
- B can verify the signature by computing
 - $V_1 = 10^{14} \bmod 19 = 16$
 - $V_2 = 4^3 \cdot 3^4 = 5184 = 16 \bmod 19$
 - Since $16 = 16$ signature is verified and valid.

Any other user can verify the signature as follows.

1. Compute $x' = a^y v^e \bmod p$.
2. Verify that $e = H(M || x')$.

To see that the verification works, observe that $x' \equiv a^y v^e \equiv a^y a^{-se} \equiv a^{y-se} \equiv a^r \equiv x \pmod{p}$

Hence, $H(M || x') = H(M || x)$