

TRACTABLE & INTRACTABLE PROBLEM

Tractable Problem: A problem that is solvable by a polynomial-time algorithm.

The upper bound is polynomial.

Here are **examples** of tractable problems (ones with known polynomial-time algorithms):

- Searching an unordered list
- Searching an ordered list
- Sorting a list
- Multiplication of integers (even though there's a gap)
- Finding a minimum spanning tree in a graph (even though there's a gap)

Intractable Problem: a problem that cannot be solved by a polynomial-time algorithm. The lower bound is exponential.

From a computational complexity stance, intractable problems are problems for which there exist no efficient algorithms to solve them.

Most intractable problems have an algorithm that provides a solution, and that algorithm is the brute-force search.

This algorithm, however, does not provide an efficient solution and is, therefore, not feasible for computation with anything more than the smallest input.

Examples

Towers of Hanoi: we can prove that any algorithm that solves this problem must have a worst-case running time that is at least $2^n - 1$.

* List all permutations (all possible orderings) of n numbers.

POLYNOMIAL-TIME ALGORITHMS: A polynomial-time algorithm is an algorithm whose execution time is either given by a polynomial on the size of the input, or can be bounded by such a polynomial. Problems which can be solved by a polynomial-time algorithm are called "tractable" problems. As an example, most algorithms on arrays can use the array size, n , as the input size. In order to find the largest element in any array requires a single pass through the array, so the algorithm which does this is of $O(n)$, or it is a "linear time" algorithm.

Sorting algorithms take $O(n \log n)$ or $O(n^2)$ time. Bubble sort takes linear time in the least case, but $O(n^2)$ time in the average and worst cases. Heapsort takes $O(n \log n)$ time in all cases. Quicksort takes $O(n \log n)$ time on average, but $O(n^2)$ time in the worst case.

As far as $O(n \log n)$ is concerned, it must be noted that the base of the logarithms is irrelevant, as the difference is a constant factor, which is ignored. All programming tasks we know have polynomial solutions. It is not due to the reason that all practical problems have polynomial-time solutions.

Rather, it is because the day-to-day problems are one for which there is no known practical solution.

NON-DETERMINISTIC POLYNOMIAL TIME ALGORITHMS: A nondeterministic computation is viewed as:

1. when a choice point is reached, an infallible oracle can be consulted to determine the right option.
2. When a choice point is reached, all choices are made and computation can proceed simultaneously.

A Non-deterministic Polynomial Time Algorithm is one that can be executed in polynomial time on a nondeterministic machine. The machine can either consult an oracle in constant time, or it can spawn an arbitrarily large number of parallel processes, which is obviously a nice machine to have.