

The Rabin-Karp-Algorithm

The Rabin-Karp string matching algorithm calculates a hash value for the pattern, as well as for each M-character subsequences of text to be compared. If the hash values are unequal, the algorithm will determine the hash value for next M-character sequence. If the hash values are equal, the algorithm will analyze the pattern and the M-character sequence. In this way, there is only one comparison per text subsequence, and character matching is only required when the hash values match.

RABIN-KARP-MATCHER (T, P, d, q)

1. $n \leftarrow \text{length}[T]$
2. $m \leftarrow \text{length}[P]$
3. $h \leftarrow d^{m-1} \bmod q$
4. $p \leftarrow 0$
5. $t_0 \leftarrow 0$
6. for $i \leftarrow 1$ to m
7. do $p \leftarrow (dp + P[i]) \bmod q$
8. $t_0 \leftarrow (dt_0 + T[i]) \bmod q$
9. for $s \leftarrow 0$ to $n-m$
10. do if $p = t_s$
11. then if $P[1\dots m] = T[s+1\dots s+m]$
12. then "Pattern occurs with shift" s
13. If $s < n-m$
14. then $t_{s+1} \leftarrow (d(t_s - T[s+1]h) + T[s+m+1]) \bmod q$

Example: For string matching, working module $q = 11$, how many spurious hits does the Rabin-Karp matcher encounters in Text $T = 31415926535\dots$

1. $T = 31415926535\dots$
2. $P = 26$
3. Here $T.Length = 11$ so $Q = 11$
4. And $P \bmod Q = 26 \bmod 11 = 4$
5. Now find the exact match of $P \bmod Q\dots$

Solution:

T =

3	1	4	1	5	9	2	6	5	3	5
---	---	---	---	---	---	---	---	---	---	---

P =

2	6
---	---

S = 0 →

3	1	4	1	5	9	2	6	5	3	5
---	---	---	---	---	---	---	---	---	---	---

$31 \bmod 11 = 9$ not equal to 4

S = 1 →

3	1	4	1	5	9	2	6	5	3	5
---	---	---	---	---	---	---	---	---	---	---

$14 \bmod 11 = 3$ not equal to 4

S = 2 →

3	1	4	1	5	9	2	6	5	3	5
---	---	---	---	---	---	---	---	---	---	---

$41 \bmod 11 = 8$ not equal to 4

S = 3 →

3	1	4	1	5	9	2	6	5	3	5
---	---	---	---	---	---	---	---	---	---	---

$15 \bmod 11 = 4$ equal to 4 SPURIOUS HIT

S = 4 →

3	1	4	1	5	9	2	6	5	3	5
---	---	---	---	---	---	---	---	---	---	---

$59 \bmod 11 = 4$ equal to 4 SPURIOUS HIT

S = 5 →

3	1	4	1	5	9	2	6	5	3	5
---	---	---	---	---	---	---	---	---	---	---

$92 \bmod 11 = 4$ equal to 4 SPURIOUS HIT

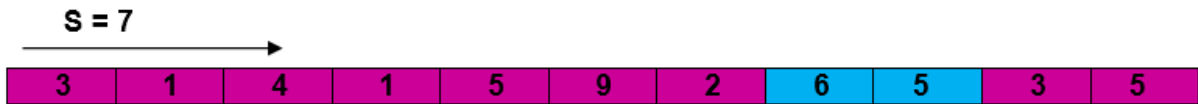
S = 6 →

3	1	4	1	5	9	2	6	5	3	5
---	---	---	---	---	---	---	---	---	---	---

$26 \bmod 11 = 4$ EXACT MATCH

S = 7 →

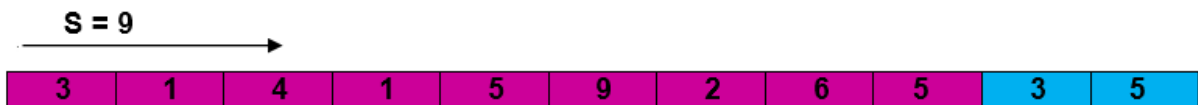
3	1	4	1	5	9	2	6	5	3	5
---	---	---	---	---	---	---	---	---	---	---



$65 \bmod 11 = 10$ not equal to 4



$53 \bmod 11 = 9$ not equal to 4



$35 \bmod 11 = 2$ not equal to 4

The Pattern occurs with shift 6.

Complexity:

The running time of **RABIN-KARP-MATCHER** in the worst case scenario $O((n-m+1)m)$ but it has a good average case running time. If the expected number of strong shifts is small $O(1)$ and prime q is chosen to be quite large, then the Rabin-Karp algorithm can be expected to run in time $O(n+m)$ plus the time to require to process spurious hits.