

SUBSET SUM PROBLEM

The *subset-sum problem* finds a subset of a given set $A = \{a_1, \dots, a_n\}$ of n positive integers whose sum is equal to a given positive integer d . For example, for $A = \{1, 2, 5, 6, 8\}$ and $d = 9$, there are two solutions: $\{1, 2, 6\}$ and $\{1, 8\}$. Of course, some instances of this problem may have no solutions.

It is convenient to sort the set's elements in increasing order. So, we will assume that $a_1 < a_2 < \dots < a_n$.

$A = \{3, 5, 6, 7\}$ and $d = 15$ of the subset-sum problem. The number inside a node is the sum of the elements already included in the subsets represented by the node. The inequality below a leaf indicates the reason for its termination.

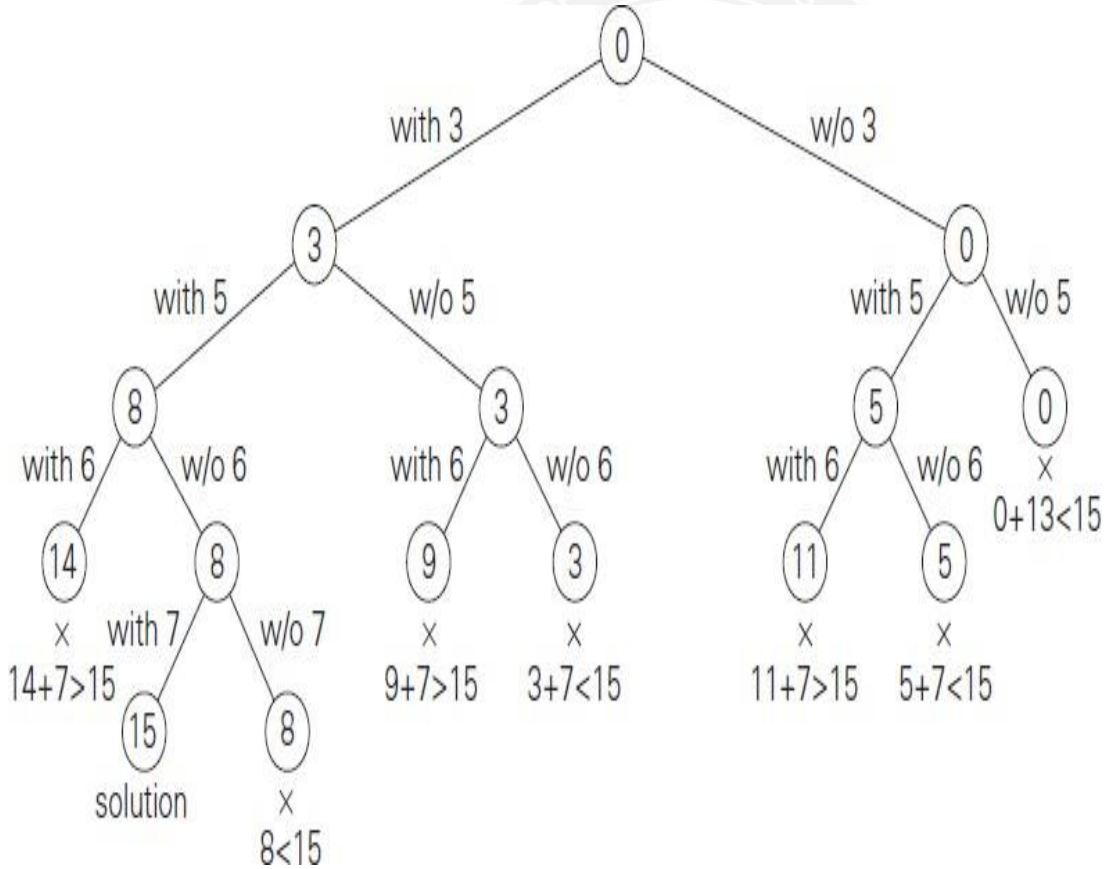


FIGURE Complete state-space tree of the backtracking algorithm applied to the instance

ALGORITHM *Backtrack*($X[1..i]$)

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//Gives a template of a generic backtracking algorithm
//Input:  $X[1..i]$  specifies first  $i$  promising components of a solution
//Output: All the tuples representing the problem's solutions
if  $X[1..i]$  is a solution write  $X[1..i]$ 

else //see Problem this section
    for each element  $x \in S_{i+1}$  consistent with  $X[1..i]$  and the constraints do
         $X[i + 1] \leftarrow$ 
         $x$ 
        Backtrack(
         $X[1..i + 1]$ )

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General Remarks

From a more general perspective, most backtracking algorithms fit the following description. An output of a backtracking algorithm can be thought of as an n -tuple (x_1, x_2, \dots, x_n) where each coordinate x_i is an element of some finite linearly ordered set S_i . For example, for the n -queens problem, each S_i is the set of integers (column numbers) 1 through n .

A backtracking algorithm generates, explicitly or implicitly, a state-space tree; its nodes represent partially constructed tuples with the first i coordinates defined by the earlier actions of the algorithm. If such a tuple (x_1, x_2, \dots, x_i) is not a solution, the algorithm finds the next element in S_{i+1} that is consistent with the values of $((x_1, x_2, \dots, x_i)$ and the problem's constraints, and adds it to the tuple as its $(i + 1)$ st coordinate. If such an element does not exist, the algorithm backtracks to consider the next value of x_i , and soon.