

SYNTAX DIRECTED DEFINITION

A syntax-directed definition specifies the values of attributes by associating semantic rules with the grammar productions. For example, an infix-to-postfix translator might have a production and rule.

PRODUCTION	SEMANTIC RULE
$E \rightarrow E_1 + T$	$E.code = E_1.code \parallel T.code \parallel '+'$

A syntax-directed translation scheme embeds program fragments called semantic actions within production bodies, as in

$$E \rightarrow E_1 + T \{ \text{print '+'} \}$$

The most general approach to syntax-directed translation is to construct a parse tree or a syntax tree, and then to compute the values of attributes at the nodes of the tree by visiting the nodes of the tree.

SDD

A syntax-directed definition (SDD) is a context-free grammar together with attributes and rules. Attributes are associated with grammar symbols and rules are associated with productions. If X is a symbol and a is one of its attributes, then we write $X.a$ to denote the value of a at a particular parse-tree node labeled X .

Inherited and Synthesized Attributes

There are two kinds of attributes for nonterminals:

Synthesized attribute:

- A synthesized attribute for a nonterminal A at a parse-tree node N is defined by a semantic rule associated with the production at N .
- Note that the production must have A as its head.
- A synthesized attribute at node N is defined only in terms of attribute values at the children of N and at N itself.

Inherited Attribute :

- An inherited attribute for a nonterminal B at a parse-tree node N is defined by

a semantic rule associated with the production at the **parent of N**.

- Note that the production must have **B** as a symbol in its body.
- An inherited attribute at node **N** is defined only in terms of attribute values at **N's parent** , **N itself**, and **N's siblings**.

Example:

PRODUCTION	SEMANTIC RULES
1) $L \rightarrow E n$	$L.val = E.val$
2) $E \rightarrow E_1 + T$	$E.val = E_1.val + T.val$
3) $E \rightarrow T$	$E.val = T.val$
4) $T \rightarrow T_1 * F$	$T.val = T_1.val \times F.val$
5) $T \rightarrow F$	$T.val = F.val$
6) $F \rightarrow (E)$	$F.val = E.val$
7) $F \rightarrow \text{digit}$	$F.val = \text{digit.lexval}$

EVALUATION ORDERS FOR SYNTAX DIRECTED DEFINITIONS

"Dependency graphs" are a useful tool for determining an evaluation order for the attribute instances in a given parse tree. While an annotated parse tree shows the values of attributes, a dependency graph helps us determine how those values can be computed. The two important classes of SDD are the "**S-attributed**" and "**L-attributed**" SDD's.

Dependency Graphs

- A dependency graph depicts the flow of information among the attribute instances in a particular parse tree; an edge from one attribute instance to another means that the value of the first is needed to compute the second.
- Edges express constraints implied by the semantic rules.
 - ✓ For each parse-tree node, say a node labeled by grammar symbol **X**, the dependency graph has a node for each attribute associated with **X**.
 - ✓ Suppose that a semantic rule associated with a production **p** defines the value of synthesized attribute **A.b** in terms of the value of **X.C**. Then, the dependency graph has an edge from **X.C** to **A.b**.
 - ✓ Suppose that a semantic rule associated with a production **p** defines the value of inherited attribute **B.c** in terms of the value of **X.a**. Then, the dependency graph has an edge from **X.a** to **B.c**.

Example:

PRODUCTION
 $E \rightarrow E_1 + T$

SEMANTIC RULE
 $E.val = E_1.val + T.val$

