

CONVERSION FROM NFA TO DFA

In NFA, when a specific input is given to the current state, the machine goes to multiple states. It can have zero, one or more than one move on a given input symbol. On the other hand, in DFA, when a specific input is given to the current state, the machine goes to only one state. DFA has only one move on a given input symbol.

Let, $M = (Q, \Sigma, \delta, q_0, F)$ is an NFA which accepts the language $L(M)$. There should be equivalent DFA denoted by $M' = (Q', \Sigma', q_0', \delta', F')$ such that $L(M) = L(M')$.

Steps for converting NFA to DFA:

Step 1: Initially $Q' = \phi$

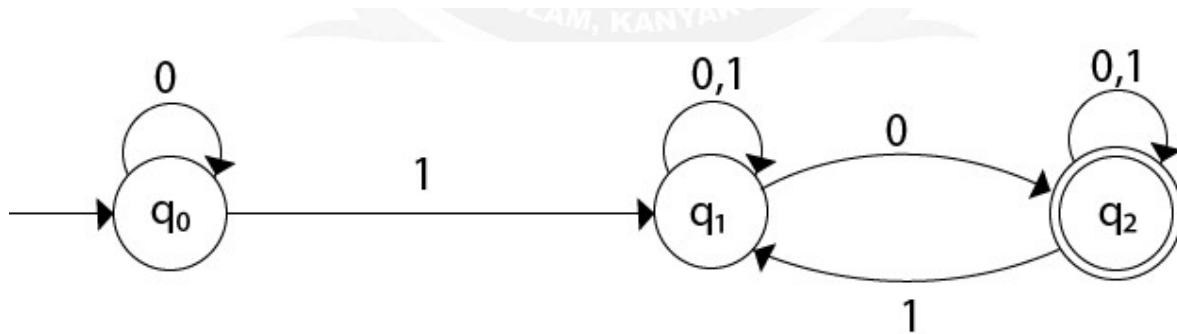
Step 2: Add q_0 of NFA to Q' . Then find the transitions from this start state.

Step 3: In Q' , find the possible set of states for each input symbol. If this set of states is not in Q' , then add it to Q' .

Step 4: In DFA, the final state will be all the states which contain F (final states of NFA)

Example 1:

Convert the given NFA to DFA.



Solution: For the given transition diagram we will first construct the transition table.

State	0	1
q0	q0	q1
q1	q1	q2
q2	q2	q1

$\rightarrow q_0$		q_0	q_1
q_1		$\{q_1, q_2\}$	q_1
$*q_2$		q_2	$\{q_1, q_2\}$

Now we will obtain δ' transition for state q_0 .

1. $\delta'([q_0], 0) = [q_0]$
2. $\delta'([q_0], 1) = [q_1]$

The δ' transition for state q_1 is obtained as:

1. $\delta'([q_1], 0) = [q_1, q_2]$ (**new state generated**)
2. $\delta'([q_1], 1) = [q_1]$

The δ' transition for state q_2 is obtained as:

1. $\delta'([q_2], 0) = [q_2]$
2. $\delta'([q_2], 1) = [q_1, q_2]$

Now we will obtain δ' transition on $[q_1, q_2]$.

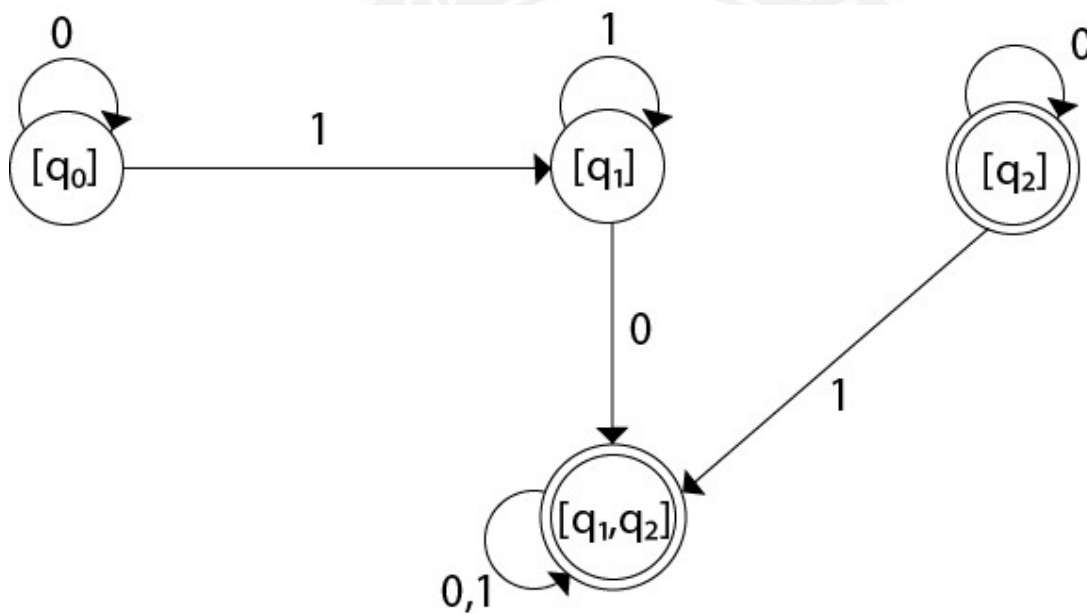
1. $\delta'([q_1, q_2], 0) = \delta(q_1, 0) \cup \delta(q_2, 0)$
2. $= \{q_1, q_2\} \cup \{q_2\}$
3. $= [q_1, q_2]$
4. $\delta'([q_1, q_2], 1) = \delta(q_1, 1) \cup \delta(q_2, 1)$
5. $= \{q_1\} \cup \{q_1, q_2\}$
6. $= \{q_1, q_2\}$
7. $= [q_1, q_2]$

The state $[q_1, q_2]$ is the final state as well because it contains a final state q_2 . The transition table for the constructed DFA will be:

State	0	1
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$\rightarrow[q_0]$	$[q_0]$	$[q_1]$
$[q_1]$	$[q_1, q_2]$	$[q_1]$
$*[q_2]$	$[q_2]$	$[q_1, q_2]$
$*[q_1, q_2]$	$[q_1, q_2]$	$[q_1, q_2]$

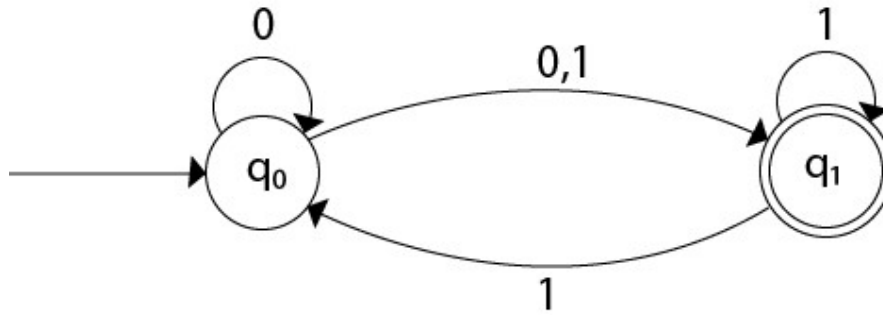
The Transition diagram will be:



The state q_2 can be eliminated because q_2 is an unreachable state.

Example 2:

Convert the given NFA to DFA.



Solution: For the given transition diagram we will first construct the transition table.

State	0	1
$\rightarrow q_0$	$\{q_0, q_1\}$	$\{q_1\}$
$*q_1$	Φ	$\{q_0, q_1\}$

Now we will obtain δ' transition for state q_0 .

- $\delta'([q_0], 0) = \{q_0, q_1\}$
- $\quad = [q_0, q_1]$ (new state generated)
- $\delta'([q_0], 1) = \{q_1\} = [q_1]$

The δ' transition for state q_1 is obtained as:

- $\delta'([q_1], 0) = \phi$
- $\delta'([q_1], 1) = [q_0, q_1]$

Now we will obtain δ' transition on $[q_0, q_1]$.

- $\delta'([q_0, q_1], 0) = \delta(q_0, 0) \cup \delta(q_1, 0)$
- $\quad = \{q_0, q_1\} \cup \phi$
- $\quad = \{q_0, q_1\}$
- $\quad = [q_0, q_1]$

Similarly,

- $\delta'([q_0, q_1], 1) = \delta(q_0, 1) \cup \delta(q_1, 1)$

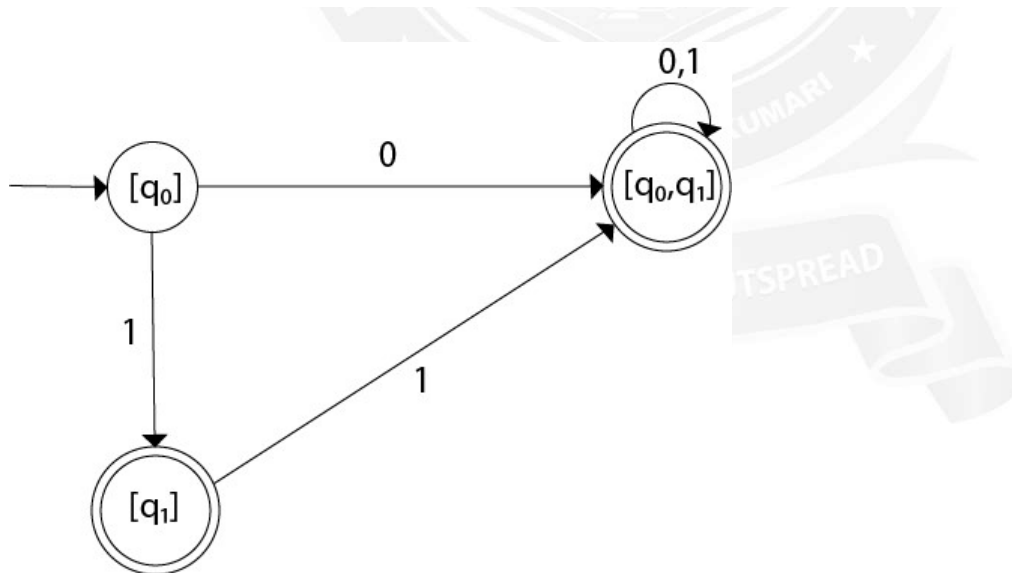
2. $= \{q1\} \cup \{q0, q1\}$
3. $= \{q0, q1\}$
4. $= [q0, q1]$

As in the given NFA, $q1$ is a final state, then in DFA wherever, $q1$ exists that state becomes a final state. Hence in the DFA, final states are $[q1]$ and $[q0, q1]$. Therefore set of final states $F = \{[q1], [q0, q1]\}$.

The transition table for the constructed DFA will be:

State	0	1
$\rightarrow[q0]$	$[q0, q1]$	$[q1]$
$*[q1]$	ϕ	$[q0, q1]$
$*[q0, q1]$	$[q0, q1]$	$[q0, q1]$

The Transition diagram will be:



Even we can change the name of the states of DFA.

Suppose

1. $A = [q_0]$
2. $B = [q_1]$
3. $C = [q_0, q_1]$

With these new names the DFA will be as follows:

