

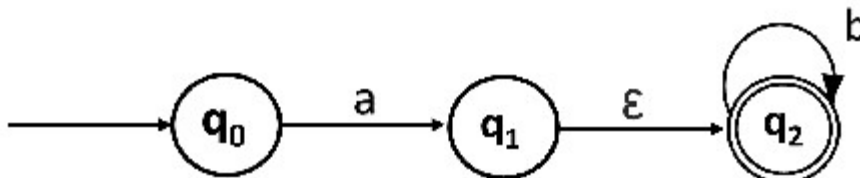
## Eliminating $\epsilon$ Transitions

NFA with  $\epsilon$  can be converted to NFA without  $\epsilon$ , and this NFA without  $\epsilon$  can be converted to DFA. To do this, we will use a method, which can remove all the  $\epsilon$  transition from given NFA. The method will be:

1. Find out all the  $\epsilon$  transitions from each state from Q. That will be called as  $\epsilon$ -closure $\{q_1\}$  where  $q_i \in Q$ .
2. Then  $\delta'$  transitions can be obtained. The  $\delta'$  transitions mean a  $\epsilon$ -closure on  $\delta$  moves.
3. Repeat Step-2 for each input symbol and each state of given NFA.
4. Using the resultant states, the transition table for equivalent NFA without  $\epsilon$  can be built.

### Example:

Convert the following NFA with  $\epsilon$  to NFA without  $\epsilon$ .



**Solutions:** We will first obtain  $\epsilon$ -closures of  $q_0$ ,  $q_1$  and  $q_2$  as follows:

1.  $\epsilon$ -closure( $q_0$ ) =  $\{q_0\}$
2.  $\epsilon$ -closure( $q_1$ ) =  $\{q_1, q_2\}$
3.  $\epsilon$ -closure( $q_2$ ) =  $\{q_2\}$

Now the  $\delta'$  transition on each input symbol is obtained as:

1.  $\delta'(q_0, a) = \epsilon$ -closure( $\delta(\delta^{\wedge}(q_0, \epsilon), a)$ )
2.  $= \epsilon$ -closure( $\delta(\epsilon$ -closure( $q_0$ ),  $a$ ))
3.  $= \epsilon$ -closure( $\delta(q_0, a)$ )
4.  $= \epsilon$ -closure( $q_1$ )
5.  $= \{q_1, q_2\}$
- 6.
7.  $\delta'(q_0, b) = \epsilon$ -closure( $\delta(\delta^{\wedge}(q_0, \epsilon), b)$ )
8.  $= \epsilon$ -closure( $\delta(\epsilon$ -closure( $q_0$ ),  $b$ ))

9.  $= \epsilon\text{-closure}(\delta(q_0, b))$
10.  $= \Phi$

Now the  $\delta'$  transition on  $q_1$  is obtained as:

1.  $\delta'(q_1, a) = \epsilon\text{-closure}(\delta(\delta^{\wedge}(q_1, \epsilon), a))$
2.  $= \epsilon\text{-closure}(\delta(\epsilon\text{-closure}(q_1), a))$
3.  $= \epsilon\text{-closure}(\delta(q_1, q_2), a)$
4.  $= \epsilon\text{-closure}(\delta(q_1, a) \cup \delta(q_2, a))$
5.  $= \epsilon\text{-closure}(\Phi \cup \Phi)$
6.  $= \Phi$
- 7.
8.  $\delta'(q_1, b) = \epsilon\text{-closure}(\delta(\delta^{\wedge}(q_1, \epsilon), b))$
9.  $= \epsilon\text{-closure}(\delta(\epsilon\text{-closure}(q_1), b))$
10.  $= \epsilon\text{-closure}(\delta(q_1, q_2), b)$
11.  $= \epsilon\text{-closure}(\delta(q_1, b) \cup \delta(q_2, b))$
12.  $= \epsilon\text{-closure}(\Phi \cup q_2)$
13.  $= \{q_2\}$

The  $\delta'$  transition on  $q_2$  is obtained as:

1.  $\delta'(q_2, a) = \epsilon\text{-closure}(\delta(\delta^{\wedge}(q_2, \epsilon), a))$
2.  $= \epsilon\text{-closure}(\delta(\epsilon\text{-closure}(q_2), a))$
3.  $= \epsilon\text{-closure}(\delta(q_2, a))$
4.  $= \epsilon\text{-closure}(\Phi)$
5.  $= \Phi$
- 6.
7.  $\delta'(q_2, b) = \epsilon\text{-closure}(\delta(\delta^{\wedge}(q_2, \epsilon), b))$
8.  $= \epsilon\text{-closure}(\delta(\epsilon\text{-closure}(q_2), b))$
9.  $= \epsilon\text{-closure}(\delta(q_2, b))$
10.  $= \epsilon\text{-closure}(q_2)$
11.  $= \{q_2\}$

Now we will summarize all the computed  $\delta'$  transitions:

1.  $\delta'(q_0, a) = \{q_0, q_1\}$

2.  $\delta'(q_0, b) = \Phi$
3.  $\delta'(q_1, a) = \Phi$
4.  $\delta'(q_1, b) = \{q_2\}$
5.  $\delta'(q_2, a) = \Phi$
6.  $\delta'(q_2, b) = \{q_2\}$

The transition table can be:

States	A	B
$\rightarrow q_0$	$\{q_1, q_2\}$	$\Phi$
* $q_1$	$\Phi$	$\{q_2\}$
* $q_2$	$\Phi$	$\{q_2\}$

**State  $q_1$  and  $q_2$  become the final state** as  $\epsilon$ -closure of  $q_1$  and  $q_2$  contain the final state  $q_2$ . The NFA can be shown by the following transition diagram:

