

SPECIFICATION OF A SIMPLE TYPECHECKER

A type checker for a simple language checks the type of each identifier. The type checker is a translation scheme that synthesizes the type of each expression from the types of its subexpressions. The type checker can handle arrays, pointers, statements and functions.

A Simple Language

Consider the following grammar:

$P \rightarrow D ; E$
 $D \rightarrow D ; D \mid id : T$
 $T \rightarrow char \mid integer \mid array [num] of T \mid \uparrow T$
 $E \rightarrow literal \mid num \mid id \mid E mod E \mid E [E] \mid E \uparrow$

Translation scheme:

$P \rightarrow D ; E$
 $D \rightarrow D ; D$
 $D \rightarrow id:T \quad \{$
 $\quad addtype(id.entry, T.type) \}$
 $T \rightarrow char \quad \{ T.type =$
 $\quad char \}$
 $T \rightarrow integer \quad \{ T.type = integer \}$
 $T \rightarrow \uparrow T1 \quad \{ T.type = pointer(T1.type) \}$
 $T \rightarrow array [num] of T1 \{ T.type := array (1... num.val , T1.type) \}$

In the above language,

- There are two basic types : char and integer ; → type_error is used to signal errors;
- the prefix operator $\uparrow$ builds a pointer type. Example , $\uparrow integer$ leads to the type expression $pointer (integer)$.

Type checking of expressions

In the following rules, the attribute type for E gives the type expression assigned to the expression generated by E.

1. $E \rightarrow literal \{$
 $\quad E.type = char \}$
 $E \rightarrow num \quad \{$
 $\quad E.type = integer \}$

Here, constants represented by the tokens literal and num have type char and integer.

2. $E \rightarrow id \quad \{ E.type = lookup (id.entry) \}$

lookup (e) is used to fetch the type saved in the symbol table entry pointed to by e.

3. $E \rightarrow E1 mod E2 \quad \{ E.type = if E1.type = integer and$
 $\quad E2.type =$
 $\quad integer then$

integer else
type_error}

The expression formed by applying the mod operator to two subexpressions of type integer has type integer; otherwise, its type is type_error.

4. $E \rightarrow E1 [E2]$ { E.type = if E2.type = integer and
E1.type
=
array(s,t)
then t
else
type_err
or }

In an array reference $E1 [E2]$, the index expression E2 must have type integer. The result is the element type t obtained from the type array(s,t) of E1.

5. $E \rightarrow E1 \uparrow$ {E.type = if E1.type =
pointer(t) then t else
type_error}

The postfix operator $\uparrow$ yields the object pointed to by its operand. The type of $E \uparrow$ is the type t of the object pointed to by the pointer E.

Type checking of statements

Statements do not have values; hence the basic type void can be assigned to them. If an error is detected within a statement, then type_error is assigned. Translation scheme for checking the type of statements:

1. Assignment statement:

$S \rightarrow id := E$ { S.type = if id.type = E.type then void
else type_error }

2. Conditional statement:

$S \rightarrow \text{if } E \text{ then } S1$ { S.type = if E.type = boolean then S1.type
else type_error }

3. While statement:

$S \rightarrow \text{while } E \text{ do } S1$ { S.type = if E.type = boolean then S1.type
else type_error }

4. Sequence of statements:

$S \rightarrow S1; S2$ { S.type = if S1.type = void and
S1.type = void then void else type_error }

Type checking of functions

The rule for checking the type of a function

application is : $E \rightarrow E1 (E2)$ { E.type:
= if E2.type = s and

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E1.type = s → t then t  
else type_error }
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