

PHYSICAL CLOCK SYNCHRONIZATION: NETWORK TIME PROTOCOL (NTP)

Centralized systems do not need clock synchronization, as they work under a common clock. But the distributed systems do not follow common clock: each system functions based on its own internal clock and its own notion of time. The time in distributed systems is measured in the following contexts:

- The time of the day at which an event happened on a specific machine in the network.
- The time interval between two events that happened on different machines in the network.
- The relative ordering of events that happened on different machines in the network.

Clock synchronization is the process of ensuring that physically distributed processors have a common notion of time.

Due to different clocks rates, the clocks at various sites may diverge with time, and periodically a clock synchronization must be performed to correct this clock skew in distributed systems. Clocks are synchronized to an accurate real-time standard like UTC (Universal Coordinated Time). Clocks that must not only be synchronized with each other but also have to adhere to physical time are termed **physical clocks**. This degree of synchronization additionally enables to coordinate and schedule actions between multiple computers connected to a common network.

Basic terminologies:

If C_a and C_b are two different clocks, then:

- **Time:** The time of a clock in a machine p is given by the function $C_p(t)$, where $C_p(t) = t$ for a perfect clock.
- **Frequency:** Frequency is the rate at which a clock progresses. The frequency at time t of clock C_a is $C_a'(t)$.
- **Offset:** Clock offset is the difference between the time reported by a clock and the real time. The offset of the clock C_a is given by $C_a(t) - t$. The offset of clock C_a relative to C_b at time $t \geq 0$ is given by $C_a(t) - C_b(t)$.

- **Skew:** The skew of a clock is the difference in the frequencies of the clock and the perfect clock. The skew of a clock C_a relative to clock C_b at time t is $C_a'(t) - C_b'(t)$.
- **Drift (rate):** The drift of clock C_a is the second derivative of the clock value with respect to time. The drift is calculated as:

$$C_a''(t) - C_b''(t).$$

Clocking Inaccuracies

Physical clocks are synchronized to an accurate real-time standard like UTC (Universal Coordinated Time). Due to the clock inaccuracy discussed above, a timer (clock) is said to be working within its specification if:

$$1 - \rho \leq \frac{dC}{dt} \leq 1 + \rho.$$

ρ - maximum skew rate.

1. Offset delay estimation

A time service for the Internet - synchronizes clients to UTC. Reliability from redundant paths, scalable, authenticates time sources. Architecture. The design of NTP involves a hierarchical tree of time servers with primary server at the root synchronizes with the UTC. The next level contains secondary servers, which act as a backup to the primary server. At the lowest level is the synchronization subnet which has the clients.

2. Clock offset and delay estimation

A source node cannot accurately estimate the local time on the target node due to varying message or network delays between the nodes. This protocol employs a very common practice of performing several trials and chooses the trial with the minimum delay.

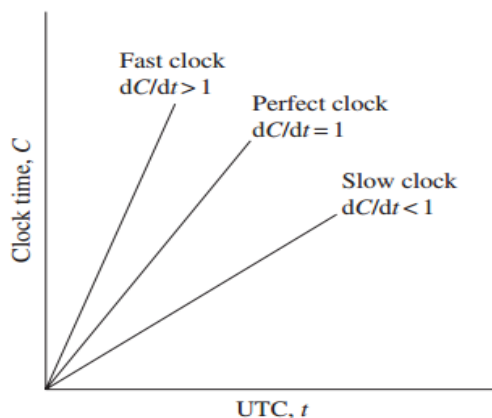


Fig : Behavior of clocks

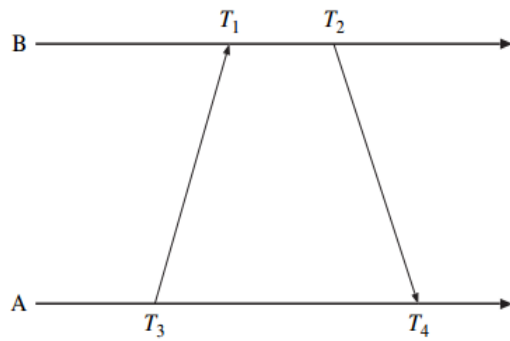


Fig a) Offset and delay estimation between processes from same server

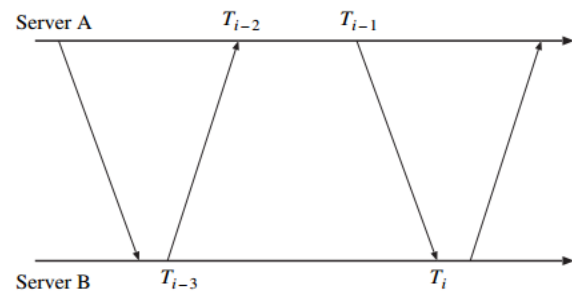


Fig b) Offset and delay estimation between processes from different servers

Let T_1, T_2, T_3, T_4 be the values of the four most recent timestamps. The clocks A and B are stable and running at the same speed. Let $a = T_1 - T_3$ and $b = T_2 - T_4$. If the network delay difference from A to B and from B to A, called **differential delay**, is small, the clock offset θ and roundtrip delay δ of B relative to A at time T_4 are approximately given by the following:

$$\theta = \frac{a+b}{2}, \quad \delta = a - b$$

Each NTP message includes the latest three timestamps T_1, T_2 , and T_3 , while T_4 is determined upon arrival.